

Software Performance Modeling of a Frame Relay Access Device

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GTE Internetworking

(work carried out at *Racal-Datacom*)

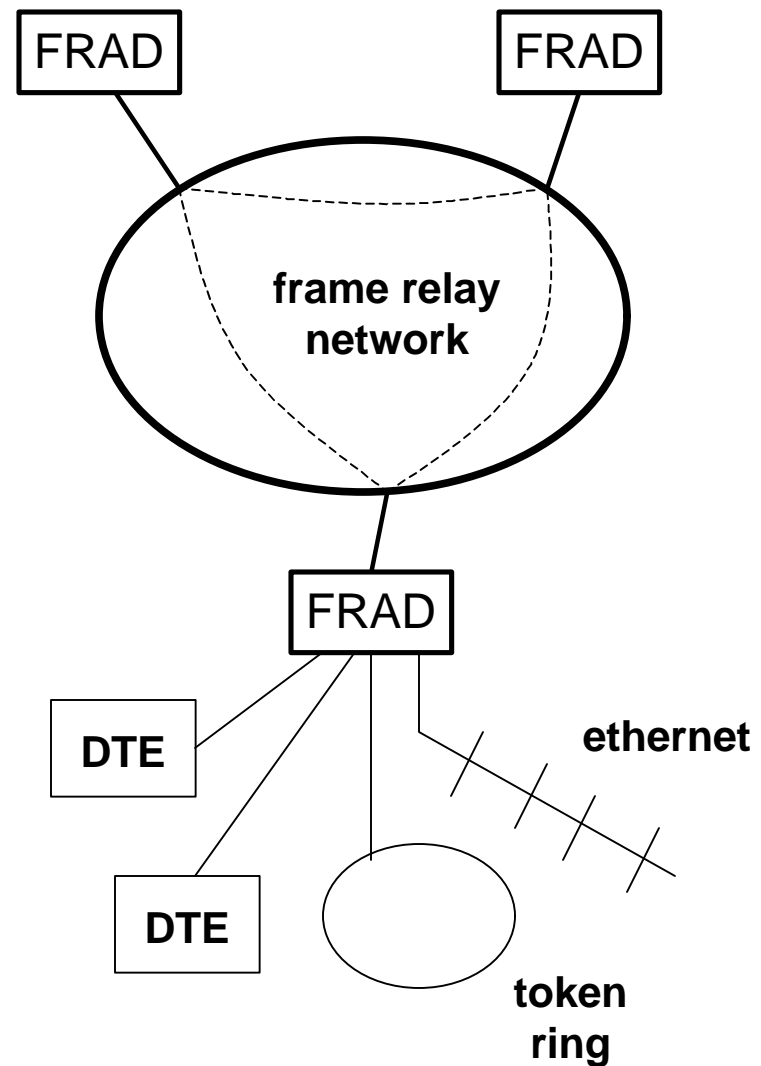
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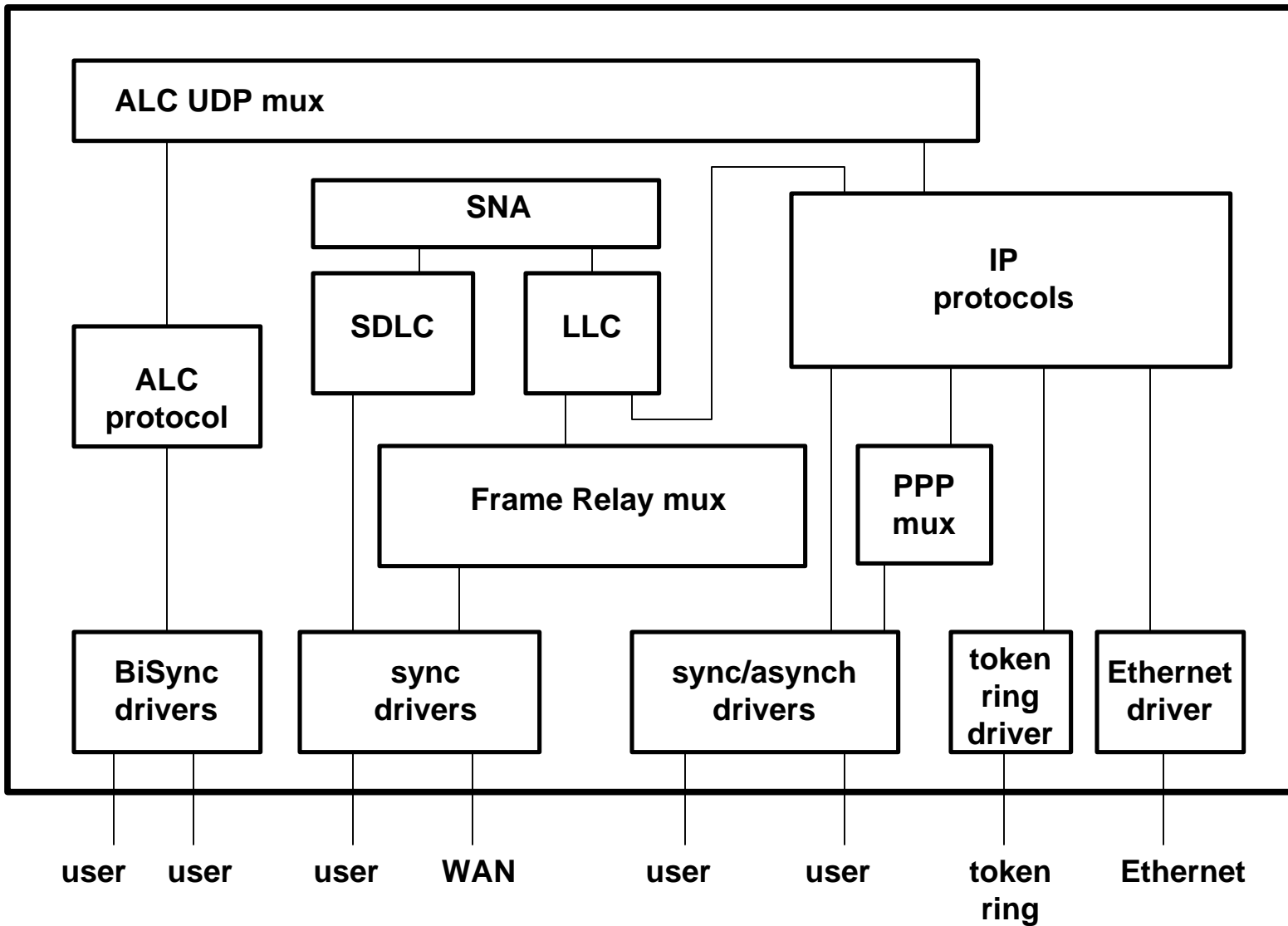
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Frame Relay Network Access Device

- **Frame Relay Network**
 - virtual-circuit service
 - connects remote sites
 - economical compared to a private leased-line network
- **FRAD**
 - interconnects LANs and DTEs to the frame relay network
 - a multi-protocol multi-function device



Racal's *FastFrame 600* FRAD Protocol Architecture



***FastFrame 600* FRAD Software**

- **Protocol Modules**
 - **standardized**
 - acquired from various vendors
 - **proprietary**
 - different software development groups

- **Protocol Integration with UNIX STREAMS Facilities**
 - **kernel routines for layered protocol software**
 - **modular architecture**
 - drivers, modules, multiplexors, queues
 - **simplifies development**
 - **reduces development time**

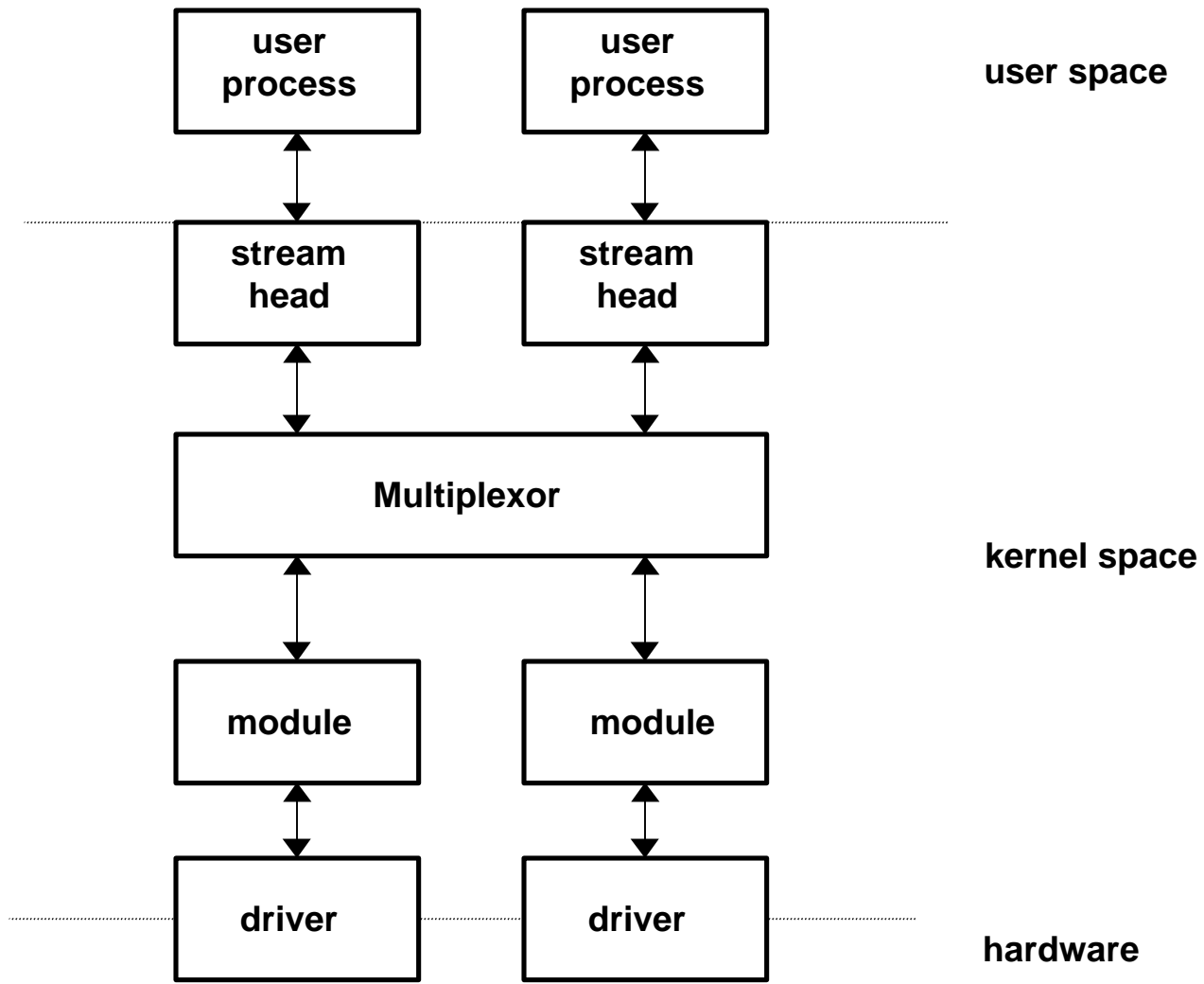
Need for Software Performance Modeling and Analysis

- **protocol modules developed by different groups**
- **performance of integrated system is unknown *a priori***
 - throughput, delay, burst handling
- **shortened time to market**
 - less time for performance measurement, re-design, tuning
- **performance ‘disasters’ at end of development cycle are costly**
- **real-time performance is of increasing importance (e.g. voice/packet)**
- **product requirements include performance**
- **need UP-FRONT analysis at product specification phase**
 - architectural choices, verify design for performance requirements, expose potential flaws
- **analysis at testing stage**
 - tool to help in parameter optimization
 - evaluate ‘last-minute’ changes

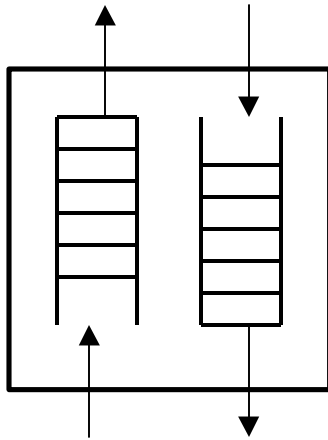
FRAD Performance Modeling and Analysis

- **We propose a software performance model for data-networking products based on STREAMS**
- **UNIX STREAMS maps naturally into a queueing model**
 - **model focuses on data-transfer phase**
 - **exploit structure imposed by STREAMS**
- **service times (path-lengths) obtained from code measurements**
- **analysis using simulation or analytical techniques**

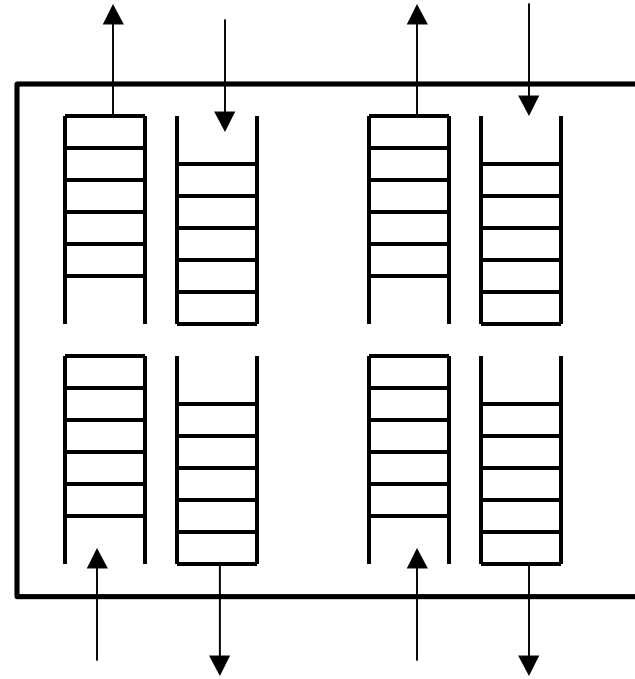
STREAMS Modeling



messages (packets)



queues in a module



queues in a multiplexor

- **Message priorities in a STREAMS queue**
 - normal messages
 - expedited messages (levels 1 to 255)
 - high-priority messages
 - FIFO scheduling within each priority band

- **Message passing from one queue to another in STREAMS**
 - involves kernel routines
 - *putnext ()*
 - *put ()*
 - *putq ()*
 - *service ()*



(1) queue A service
calls putnext

(2) putnext passes
message to queue B put

(3) queue B put
processes message

(4) put passes
message to putq

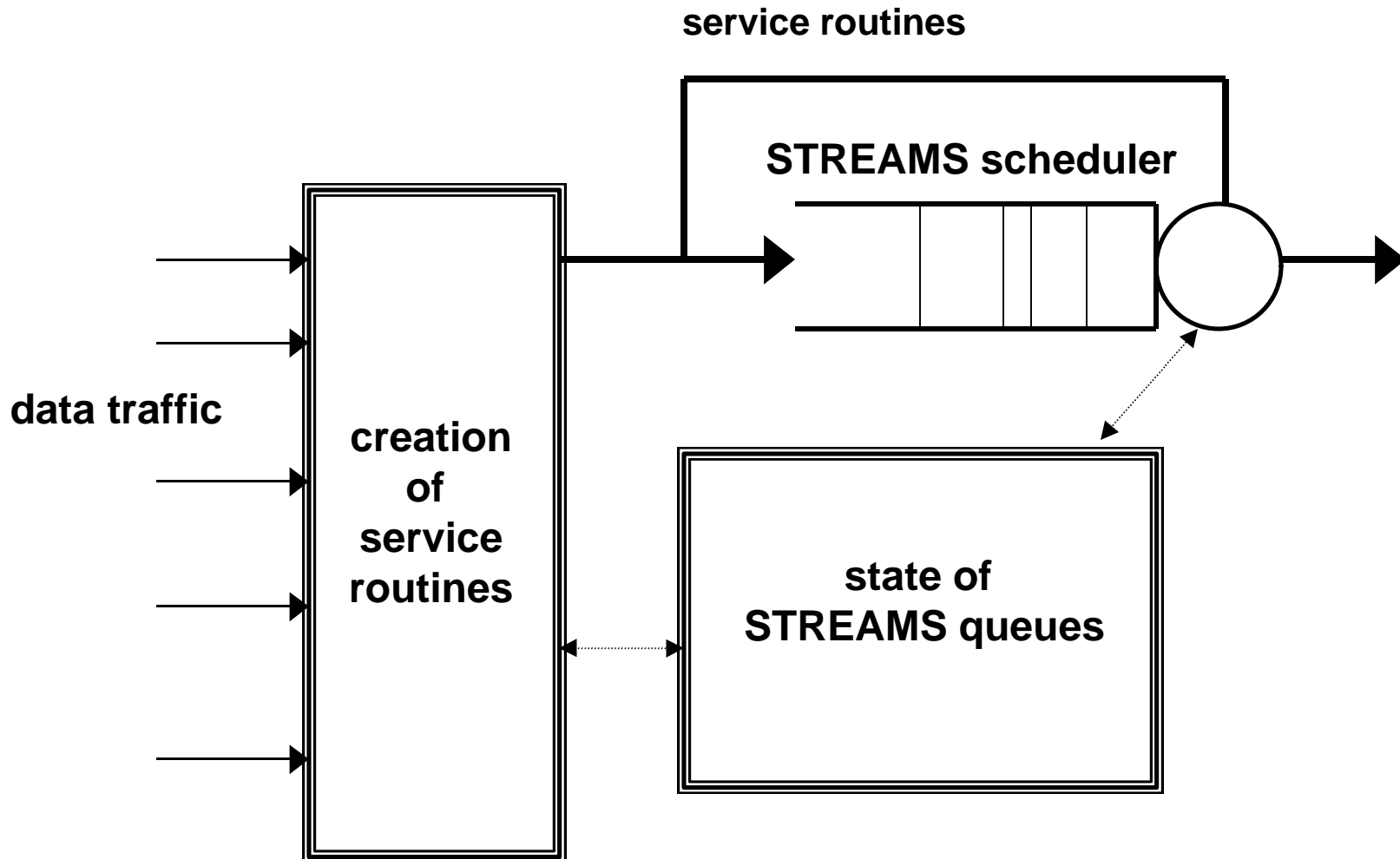
(5) putq places message
on queue B

(6) putq schedules
queue B service

- **Scheduling of Service Routines by STREAMS**
 - service routines called by STREAMS scheduler
 - STREAMS scheduler is FIFO
 - STREAMS scheduler processes all messages on a queue when service routine is called

- **Inter-Queue Flow Control in STREAMS**
 - counter *q_count*
 - high and low water marks
 - service routines “put to sleep” if flow control in force
 - service routines “woken up” when flow control removed

State-Dependent FIFO Queueing Model



Model Analysis

- **Analytical**
 - difficult due to complex state-dependencies
 - possible to develop a simplified Markov chain model
 - a challenging performance analysis problem
- **Simulation**
 - simulation model implemented using OPNET Modeler
 - could automate building of OPNET model using OPNET EMA interface

Concluding Remarks

- **Software performance modeling and analysis is an essential for data-networking product development**
- **research needed for SPE of data-networking products**
- **automated tools are needed for SPE**
- **we proposed a queueing model for SPE of STREAMS-based data-networking products**