SYSC 5701 Operating System Methods for Real-Time Applications

**Memory Issues** 

Winter 2014

# Memory Management is a Gnarly Issue

- Pearce suggests: offload the problem to the application ... how do some "real" RTOS's cope?
- µC/OS II: has (optional) API to manage partitions
  - Application supplies memory for partitions
  - Each partition contains blocks of same size
  - Partitions are managed by kernel (safe)
- FreeRTOS: has required API
  - add-on above kernel
  - various implementations



# µC/OS II

- Partition a contiguous memory region
  - Kernel cuts up into blocks of equal size
  - Application can Get/Put blocks
- Can have multiple partitions of different sized blocks
- Configuration (OS\_CFG.H)
  - Enable management services: OS\_MEM\_EN = 1
  - Max. number of partitions:
    - OS\_MAX\_MEM\_PART = MaxNumberOfPartitions
    - Statically allocated array of memory control blocks (not partitions!)

Application specifies these at config! (compile time)



# µC/OS II

Application supplies these!

Create a partition at runtime:
 OS\_MEM \*OSMemCreate(

void \*addr, // start address of memory block
INT32U nblks, // number of blocks to create
INT32U blksize, // size of each block
INT8U \*perr // return code (for create status)

• Returns: OS\_MEM\* = ptr to memory control block

);





- Get a block from a partition at runtime:
- void \*OSMemGet(
  - OS\_MEM \*pmem, // ptr to memory control block INT8U \*perr // return code (success?) // size of block is implicit to partition!!
- );
  - Returns: void\* = pointer to block from specified partition





• Put (return) a block to a partition: INT8U OSMemPut(

void \*pblk

OS\_MEM \*pmem, // ptr to memory con lock // ptr to block to return

);

Returns: INT8U = return code 



- Memory allocation API is in the portable layer
- Portable layer: outside of source files that implement the core RTOS functionality
- Allows an application-specific implementation appropriate to real time system being developed
- Provides some implementations ... but application can supply its own implementation



- FreeRTOSConfig.h customizes the kernel to the application being built
- Every FreeRTOS application must have a FreeRTOSConfig.h header file in the application directory (not RTOS directory!)
- configTOTAL\_HEAP\_SIZE = xxxx
  - Total amount of RAM available to RTOS kernel
     only used if application uses particular
     provided sample memory allocation schemes

Application specifies this at config! (compile time)



 When kernel requires RAM, calls: void \*pvPortMalloc( size\_t xWantedSize )

- Returns ptr to allocated block of requested size
- Returns NULL if no memory allocated
- When kernel releases RAM, calls: void pvPortFree( void \*pv )

• Uses whatever implementation has been linked to the kernel code



- Provided implementation: Heap\_1.c
- Does not permit memory to be freed once it has been allocated (i.e. no pvPortFree calls)
  - Deterministic
  - OK when all kernel-managed objects are created initially at startup and exist for entire running of application
  - Pros: Simple, no runtime overhead after startup
  - Cons: no dynamic create/delete



- Provided implementation: Heap\_2.c
- Best fit algorithm, allows blocks to be freed, does not coalesce adjacent free blocks to create larger blocks
  - NOT Deterministic
  - OK when kernel-managed objects are created (deleted) dynamically, but only a small set of sizes of blocks involved – e.g. fixed sized control blocks & messages
  - Pros: dynamic create/delete
  - Cons: runtime overhead, random-sized blocks will likely increase fragmentation



- Provided implementation: Heap\_3.c
- Simple thread-safe wrapper on C's malloc() & free()
  - NOT Deterministic
  - OK when kernel-managed objects are created (deleted) dynamically, and random-sized blocks
  - Must now include C library for implementation of malloc() and free() [code size & efficiency?!]
  - configTOTAL\_HEAP\_SIZE not used
  - Pros: dynamic create/delete
  - Cons: runtime overhead

- Provided implementation: Heap\_4.c
- First fit algorithm, coalescence algorithm
  - NOT Deterministic
  - OK when kernel-managed objects are created (deleted) dynamically, and random-sized blocks
  - Probably more efficient than C library (smaller code?)
  - Pros: dynamic create/delete
  - Cons: runtime overhead



- Application provided implementation: Heap\_x.c
- Implemented however the application might like to manage memory ... must implement: void \*pvPortMalloc( size\_t xWantedSize )
  - void pvPortFree( void \*pv )
  - Used by kernel when needed
  - Pros: use system-specific memory regions in customized ways (?)
  - Cons: more code/details for application programmer

